

Exercise Session 2:

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Exercise 1: Find a phrase-structure grammar for each of these languages:

1. The set consisting of the bit strings 0, 1, and 11.

Solution: We need to find a grammar G such that $L(G) = \{0, 1, 11\}$,
This language is finite so only need the starting values

$$S \rightarrow 0$$

$$S \rightarrow 1$$

$$S \rightarrow 11$$

2. All sets of bit strings containing only 1s.

Solution: We need to find a grammar G such that $L(G) = \{1, 11, \dots, \underbrace{11\dots1}_n\}$.

$$S \rightarrow 1$$

$$S \rightarrow 1S$$

3. All sets of bit strings that start with 0 and end with 1.

Solution: We need to find the grammar G such that $L(G) = \{01, 001, 011, \dots\}$

$$S \rightarrow 0A$$

$$A \rightarrow 0A$$

$$A \rightarrow 1A$$

$$A \rightarrow \lambda$$

4. All sets of bit strings that starts with a 0 followed by an even number of 1s

Solution: We need to find the grammar G such that $L(G) = \{011, 01111, \dots, 0\underbrace{11\dots1}_{2n}\}$

$$S \rightarrow 0A$$

$$A \rightarrow 11A$$

$$A \rightarrow \lambda$$

Exercise 2: Express each of these sets using a regular expression.

1. All sets consisting of strings 0, 11, and 010.

Solution: A string of this set is either 0, 11 or 010 so we simply write $0|11|010$.

2. All sets of bit strings of three 0s followed by two or more 0s

Solution: To be clear, this is the same as saying that we need at least five 00000, to represent the uncertain amount of 0s that follow we need to end the regular expression with 0^* . So in conclusion the regular expression is 000000^* .

3. All sets of bit strings of odd length

Solution: An odd number can be written as $2k + 1$ for integers k . So the first bit is represented by $0|1$. This needs to be followed by uncertain amount of even strings. Thus the regular expression would be $(0|1)((0|1)(0|1))^*$

4. All sets of bit strings that contain exactly one 1.

Solution: The string can start or end with an uncertain amount of 0s as long as there is exactly one 1, therefore the regular expression is 0^*10^*

5. Empty bit string and all sets of bit strings ending in 1 and not containing 000.

Solution: So the string cannot end with 0001, but 01 or 001 or you can have several 1s in a row. The regular expression would be $(1|01|001)^*$

Exercise 3: Let $V = \{S, A, B, a, b\}$ and $T = \{a, b\}$. Find the language generated by the grammar (V, T, S, P) when the set of P of productions consists of:

1. $S \rightarrow AB, A \rightarrow ab, B \rightarrow bb$.

Solution:

$$S \rightarrow AB \rightarrow \underbrace{ab}_A \underbrace{bb}_B = abbb$$

Thus $L(G) = \{abbb\}$

2. $S \rightarrow AB, S \rightarrow aA, A \rightarrow a, B \rightarrow ba$

Solution:

$$S \rightarrow aA \rightarrow a \underbrace{a}_A = aa$$

$$S \rightarrow AB \rightarrow \underbrace{a}_A \underbrace{ba}_B = aba$$

Thus $L(G) = \{aa, aba\}$

3. $S \rightarrow AB, S \rightarrow AA, A \rightarrow aB, A \rightarrow ab, B \rightarrow b$

Solution:

$$S \rightarrow AB \rightarrow \underbrace{A}_{aB} B \rightarrow a \underbrace{B}_b \underbrace{B}_b \rightarrow abb$$

$$S \rightarrow AA \rightarrow \underbrace{A}_{aB} \underbrace{A}_{aB} \rightarrow a \underbrace{B}_b a \underbrace{B}_b = abab$$

Exercise 4: Explain what the productions are in a grammar if the "Backus-Naur form" for productions is as follows:

$$\begin{aligned} \text{expression} &::= (\text{expression})| \\ &\quad \text{expression} + \text{expression}| \\ &\quad \text{variable} \\ \text{variable} &::= x|y \end{aligned}$$

Exercise 5: Draw a diagram belonging to the table below. The starting state

Table 1:

State	f		g	
	Input		Input	
	0	1	0	1
s_0	s_0	s_4	1	1
s_1	s_0	s_3	0	1
s_2	s_0	s_2	0	0
s_3	s_1	s_1	1	1
s_4	s_1	s_0	1	0

is state s_0 . what is the output for "1110010101"?

Solution: We start at s_0

$In = 1, s_0 \rightarrow s_4, Out = 1$

$In = 1, s_4 \rightarrow s_0, Out = 0$

$In = 1, s_0 \rightarrow s_4, Out = 1$

$In = 0, s_4 \rightarrow s_1, Out = 1$

$In = 0, s_1 \rightarrow s_0, Out = 0$

$In = 1, s_0 \rightarrow s_4, Out = 1$

$In = 0, s_4 \rightarrow s_1, Out = 1$

$In = 1, s_1 \rightarrow s_3, Out = 1$

$In = 0, s_3 \rightarrow s_1, Out = 1$

$In = 1, s_1 \rightarrow s_3, Out = 1$

So the output sequence is 1011011111

Exercise 6: Construct a deterministic finite state automation that recognizes the set of all bit strings beginning with 01

Solution: We need to find M with $L(M) = \{01x\}$ where x is any string. First we start out with and initial state s_0 . Define s_1 so that $s_0 \xrightarrow{0} s_1$, then a final state s_2 with $s_1 \xrightarrow{1} s_2$. Finally we need a state s_3 with $s_3 \xrightarrow{0,1} s_3, s_0 \xrightarrow{1} s_3$ and $s_1 \xrightarrow{0} s_3$.

Exercise 7: Show that there is no finite-state automation with two states that recognizes the set of all bit strings that have one or more 1s bits and end with a 0.

Solution: Assume there exists such a machine with $L(M) = \{10, 110, 000110, \dots\}$ and start state s_0 and another state s_1 . because the empty string is not in the language but some strings are accepted, we must have s_1 as the only final state, with at least one transition from s_0 to s_1 . Because the string 0 is not in the language, the transition from s_0 to s_1 . Because the string 0 is not in the language, the transition from s_0 on input 0 must be to itself, so the transition from s_0 on the input 1 must be to s_1 . But this contradicts the fact that 1 is not in the language.

Exercise 8: Construct a state-machine representing a unit-delay machine, which produces as output the input string delayed by a 0, that is, it produces as output bit string $0x_1x_2...x_{k-1}$ given the input bit string $x_1x_2...x_k$?

Solution:

Table 2:

State	f		g	
	Input		Input	
	0	1	0	1
s_0	s_2	s_1	0	0
s_1	s_1	s_2	1	1
s_2	s_2	s_2	0	0