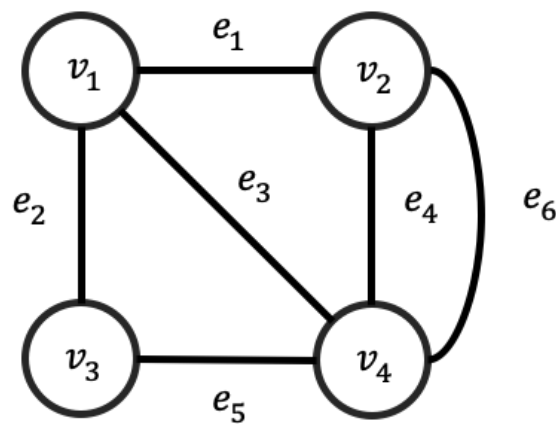


**Exercise:** List the degree of each vertex in the following graphs.

Verify  $\sum_{\{v \in V\}} \deg(v) = 2m$

**Solution:**



We have  $m = |E| = 6$ , hence  $2m = 12$ .

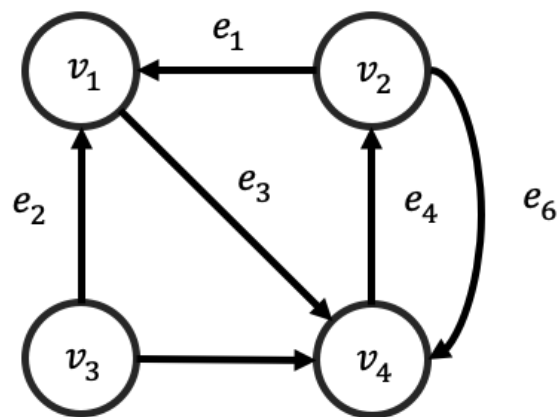
$$\begin{aligned} \sum_{\{v \in V\}} \deg(v) &= \deg(v_1) + \deg(v_2) + \deg(v_3) + \deg(v_4) \\ &= 3 + 3 + 2 + 4 = 12 \end{aligned}$$

**Exercise:** List the degree of each vertex in the following graphs. Verify

$$\sum_{\{v \in V\}} \deg^-(v) = \sum_{\{v \in V\}} \deg^+(v) = m,$$

*In degree of a vertex  $v$  ( $\deg^-(v)$ ) is the number of edges with  $v$  as their terminal vertex.*

**Definition:** *Out degree of a vertex  $v$  ( $\deg^+(v)$ ) is the number edges with  $v$  as their initial vertex.*



**Solution:** We have  $m = |E| = 6$  edges

$$\begin{aligned} \sum_{\{v \in V\}} \deg^-(v) &= \deg^-(v_1) + \deg^-(v_2) + \deg^-(v_3) + \deg^-(v_4) \\ &= 1 + 2 + 2 + 1 = 6 \end{aligned}$$

$$\begin{aligned} \sum_{\{v \in V\}} \deg^+(v) &= \deg^+(v_1) + \deg^+(v_2) + \deg^+(v_3) + \deg^+(v_4) \\ &= 2 + 1 + 0 + 3 = 6 \end{aligned}$$

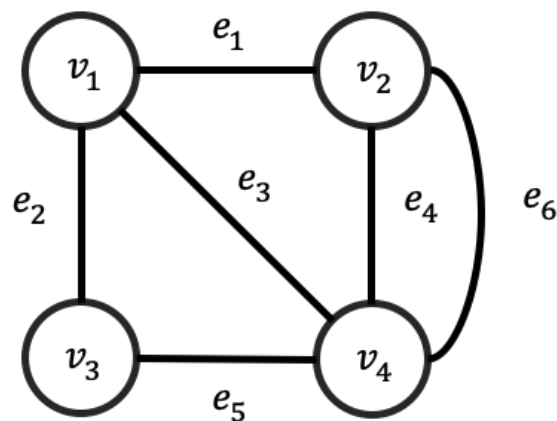
So verified!

**Exercise:** How many edges are there in an undirected graph with 10 vertices each of degree ?

**Solution:**

$$\begin{aligned}
 2m &= \sum_{\{v \in V\}} \deg(v) \\
 &= 10 \cdot 6 \\
 &= 60 \\
 \Rightarrow m &= 30
 \end{aligned}$$

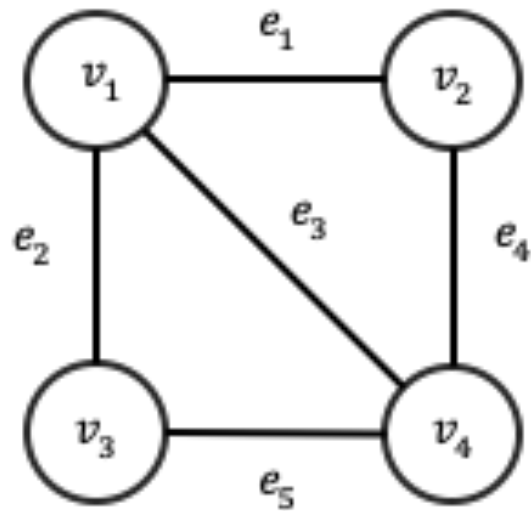
**Exercise:** Write down the adjacency matrix of the graph  $G$ .



**Solution:** Let  $n = |V|$  be the number of vertices. An adjacency matrix  $A$  of  $G$  with is an  $n \times n$  matrix with elements  $a_{ij} = 1$  if  $\{v_i, v_j\}$  is an edge of  $G$ . In this example every combination of the vertices are edges in  $G$  except for  $\{v_2, v_3\}$ , thus  $a_{23} = a_{32} = 0$ , thus

$$A = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

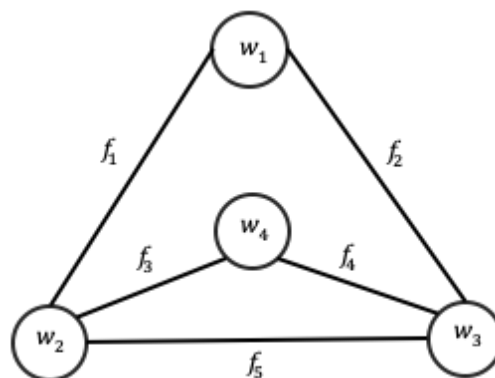
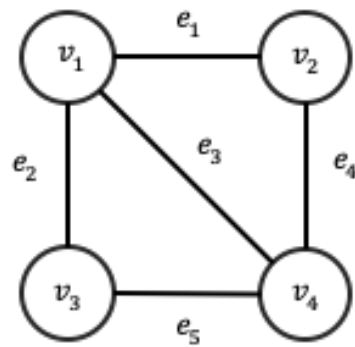
**Exercise:** Write down the incidence matrix of the graph  $G$



The incidence matrix with respect to this ordering of  $V$  and  $E$  is the  $n \times m$  matrix  $M = [m_{ij}]$  where  $m_{ij} = 1$  if edge  $e_j$  is incident with  $v_i$  thus we have

$$M = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 1 \end{bmatrix}$$

**Exercise:** Prove that the following two graphs are isomorphic



**Definition:**

The simple graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ , are said to be isomorphic if there exists a one-to-one and onto function  $f$  from  $V_1$  to  $V_2$  with the property that  $a$  and  $b$  are adjacent in  $G_1$  if and only if  $f(a)$  and  $f(b)$  are adjacent in  $G_2$ , for all  $a$  and  $b$  in  $V_1$ . Such a function  $f$  is called an isomorphism.

Two simple graphs that are not isomorphic is called nonisomorphic.

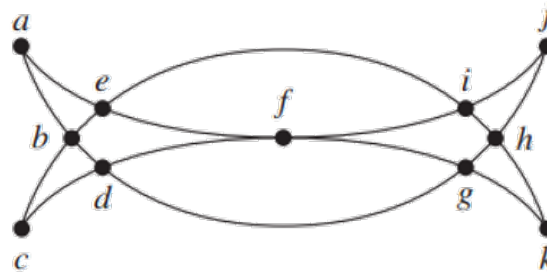
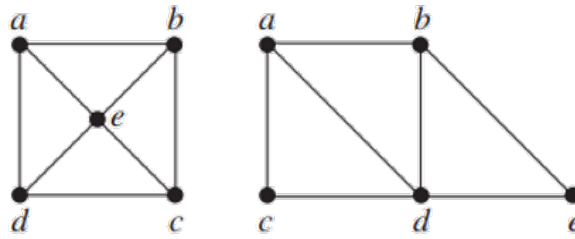
Both graphs contain four vertices and five edges. Two vertices have  $\deg = 2$  and two vertices have  $\deg = 3$  in both graphs. It is also easy to see that all subgraphs of are isomorphic. Because they agree with respect to these invariants, it is reasonable to try find an isomorphism  $f$ .

Because  $v_1$  and  $v_4$  are connected to two vertices with  $\deg = 2$  and  $v_2$  and  $v_3$  are connected to two vertices of  $\deg = 3$  we must map these to vertices with the same properties. So  $f(v_1) = w_2$  and  $f(v_4) = w_3$  as  $v_1, v_4$  and  $w_2, w_3$  have the same properties. Thus  $f(v_2) = w_1$ ,  $f(v_3) = w_4$

To see whether  $f$  preserves edges, we examine the adjacency matrices.

$$\begin{bmatrix} v_1 & v_2 & v_3 & v_4 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} \begin{bmatrix} w_2 & w_1 & w_4 & w_3 \\ 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

**Exercise:** Do the following graph contain an Euler path respectively? What about an Euler circuit?



Reminde yourself

**Definition:**

*An Euler circuit is a graph  $G$  that is a simple circuit containing every edge of  $G$ .*

*An Euler*

*path in  $G$  is a simple path containing every edge of  $G$*

Nessesary conditions for existence of Euler circuit/path:

**Theorem:**

*A connected multigraph with atleast two vertices has an Euler circuit if and only if each of its vertices has an even deegree*

*A connected multigraph has an Euler path but not an Euler circuit if and only if it has two vertices of odd degree.*

The first (upper left corner) graph does not contain an Euler circuit since four out of five vertex have odd degree. neither does it contain a Euler path.

The second (upper right corner) graph has an Euler circuit since  $\deg(a) = \deg(b) = 3$  and  $\deg(c) = 2$ ,  $\deg(d) = 4$ ,  $\deg(e) = 2$  thus we have exactly two verticies with odd degree, therefore it has an Euler path but not an Euler circuit.

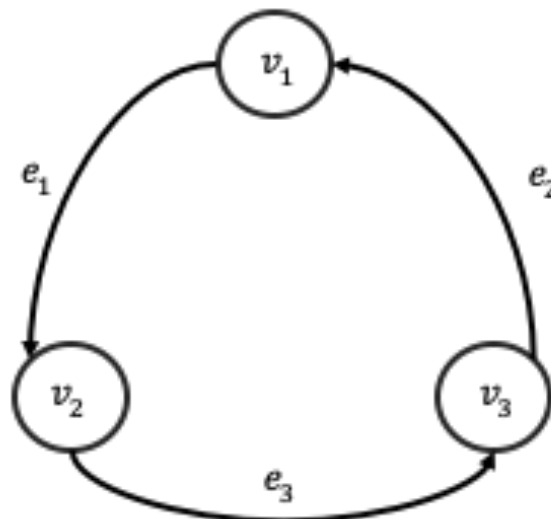
In Mohammad Scimitars each vertex has  $\deg = 2$  so it is a Euler circuit. The graph however does not have any vertex of odd degree so there is no Euler path.

**Exercise:** Is this graph strongly / weakly connected?

**Definition:**

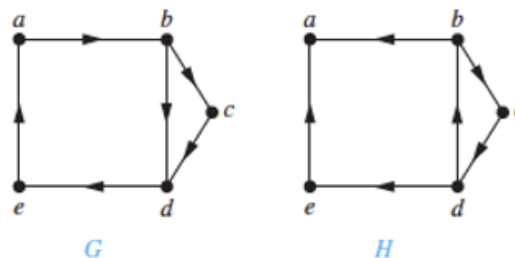
*A graph is called strongly connected if there is a path from  $a$  to  $b$  and from  $b$  to  $a$  for all vertices  $\{a, b\}$*

*A graph is called weakly connected if there is a path between every two vertices in the underlying undirected graph.*



**Solution:**

It is clearly also strongly connected because we can go from  $v_1$  to  $v_3$  by the road  $v_1 \rightarrow v_2 \rightarrow v_3$  and we can go from  $v_3$  to  $v_1$  by  $v_3 \rightarrow v_1$





$G$  is clearly strongly connected and therefore also weakly connected

$H$  is weak. If undirected it is clearly connected! If it is directed then you can go  $b \rightarrow c \rightarrow d \rightarrow e$  but there is no way to go from  $e$  to  $b$  as the only walk you can do from  $e$  is  $e \rightarrow a$ .

